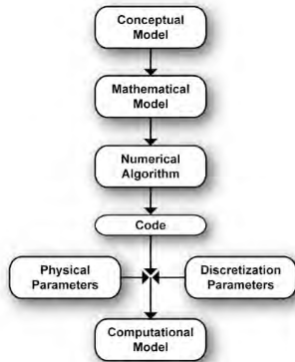


Algorithms and Operation Counting

ASEN 3502 — Computational Methods



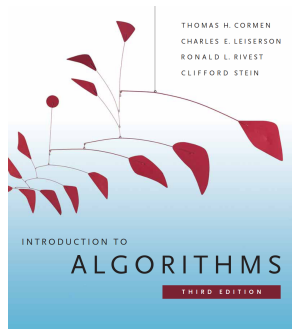
NM-11050-33

ASME V&V 10, *Guide for Verification and Validation in Computational Solid Mechanics*.

What is an algorithm?

“Informally, an algorithm is any well-defined computational **procedure** that takes some value, or set of values, as **input** and produces some value, or set of values, as **output**.”

— Cormen et al.



Introduction to Algorithms

Cormen, Leiserson,
Rivest, and Stein

mitpress.mit.edu

Another type of algorithm



Why algorithms?

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“An algorithm is said to be **correct** if, for every input instance, it halts with the correct output.” — Cormen et al.
 - ▶ We can make precise and rigorous statements without having to deal with issues like roundoff error, memory allocation, etc.

How do we represent algorithms?

Matrix-vector multiplication: three ways

English

To get entry r of $\{b\}$, multiply each entry in row r of $[A]$ by the matching entry of $\{x\}$, and add up the results. Do this for every row.

Code

```
b = zeros(n,1);  
for r = 1:n  
    for c = 1:n  
        b(r) = b(r) ...  
            + A(r,c)*x(c);  
    end  
end
```

Pseudocode

```
input      :  $n \times n$  matrix  $[A]$ ,  
              vector  $\{x\}$   
output    :  $\{b\} = [A]\{x\}$   
for  $r \leftarrow 1$  to  $n$  do  
     $b[r] \leftarrow 0$   
    for  $c \leftarrow 1$  to  $n$  do  
         $b[r] \leftarrow b[r] + A[r, c] \cdot x[c]$   
    end  
end
```

“What separates pseudocode from “real” code is that in pseudocode, we employ whatever expressive method is most clear and concise to specify a given algorithm.” — Cormen et al.

Analysis of algorithms

- ▶ *Analyzing* an algorithm means predicting the resources that an algorithm requires
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- ▶ *Analyzing* an algorithm means predicting the resources that an algorithm requires
 - ▶ For today, we will count floating point operations (“Flops”): multiplication, division, addition, and subtraction of floating point numbers
- ▶ **Discuss with your neighbor:** Consider the equation $[A]\{x\} = \{b\}$. As n gets larger, which requires more floating point operations? Finding $\{x\}$ given $[A]$ and $\{b\}$, or finding $\{b\}$ given $[A]$ and $\{x\}$?

Matrix-vector multiplication

input : $n \times n$ matrix $[A]$, vector $\{x\}$

output : $\{b\} = [A]\{x\}$

for $r \leftarrow 1$ **to** n **do**

$b[r] \leftarrow 0$

for $c \leftarrow 1$ **to** n **do**

$b[r] \leftarrow b[r] + A[r, c] \cdot x[c]$

end

end

Gaussian elimination

input : $n \times n$ matrix $[A]$, right-hand side $\{b\}$

output : solution $\{x\}$ of $[A]\{x\} = \{b\}$

// forward elimination

for $c \leftarrow 1$ **to** $n - 1$ **do**

for $r \leftarrow c + 1$ **to** n **do**

$f \leftarrow A[r, c] / A[c, c]$

for $j \leftarrow c$ **to** n **do**

$A[r, j] \leftarrow A[r, j] - f \cdot A[c, j]$

end

$b[r] \leftarrow b[r] - f \cdot b[c]$

end

end

// back substitution

for $r \leftarrow n, n - 1, \dots, 1$ **do**

$s \leftarrow b[r]$

for $j \leftarrow r + 1$ **to** n **do**

$s \leftarrow s - A[r, j] \cdot x[j]$

end

$x[r] \leftarrow s / A[r, r]$

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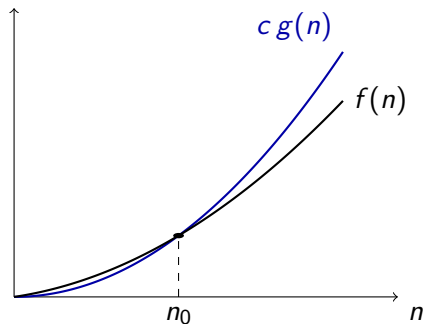
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Definition

$f(n) = O(g(n))$ as $n \rightarrow \infty$ if there exist positive constants c and n_0 such that

$$0 \leq f(n) \leq c g(n) \quad \text{for all } n \geq n_0$$



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